

Half-Double: Hammering From the Next Row Over

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- **Rowhammer**

- Default refresh window of 64 ms



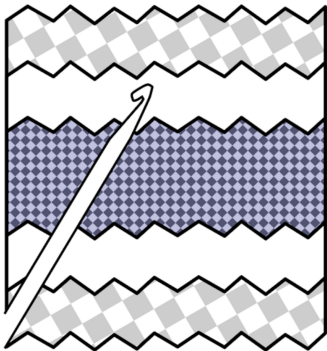
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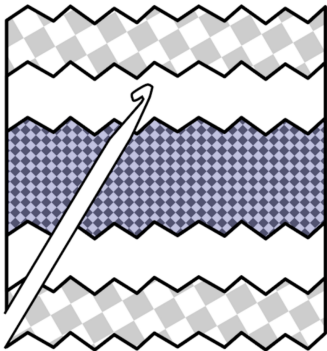
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 - Refresh direct neighbours hammering rows
 - **Exhaustion** with multi-sided patterns [2, 1]



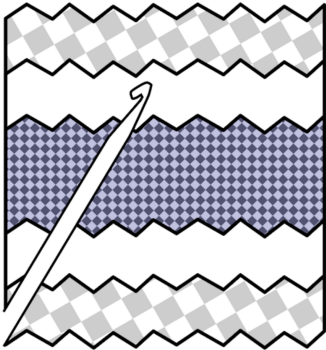
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- **Targeted Row Refresh (TRR)**
 - Refresh direct neighbours hammering rows
 - **Exhaustion** with multi-sided patterns [2, 1]
- Would perfect TRR fix Rowhammer attacks?



- Short answer: No



- Short answer: **No**
- Hammering with three rows between the aggressors
 - Causes **flips** on LPDDR4x commodity devices
 - 5 out of 7 mobile devices affected
 - With active **TRR** and on-chip **ECC**



- Short answer: **No**
- Hammering with three rows between the aggressors
 - Causes **flips** on LPDDR4x commodity devices
 - 5 out of 7 mobile devices affected
 - With active **TRR** and on-chip **ECC**
- Is this *Distance-2* Rowhammer?
- What is the **root cause**?

Far Aggressor	(\mathcal{F}_+)
Near Aggressor	(\mathcal{N}_+)
Victim	(\mathcal{V})
Near Aggressor	(\mathcal{N}_-)
Far Aggressor	(\mathcal{F}_-)

- FPGA setup
 - Control DIMM via FPGA
 - Full control over the refreshes
 - Deactivated TRR
 - No need for data retention



Far Aggressor	(\mathcal{F}_+)
Near Aggressor	(\mathcal{N}_+)
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- **Distance-1**

- $(\mathcal{N}_+ \rightarrow \mathcal{N}_-)^{\infty}$
- *Classic* double-sided Rowhammer

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- **First flip after:**
 - 18 000 hammers in 1.2 *ms*

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 - 18 000 hammers in 1.2 *ms*
 - ✓ **Within** the refresh window
 - ✗ **Mitigated** by TRR



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- **Distance-2**

- $(\mathcal{F}_+ \rightarrow \mathcal{F}_-)^{\infty}$
- Distance two double-sided Rowhammer

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 - ✗ **Not** within the refresh windows



- **Half-Double**

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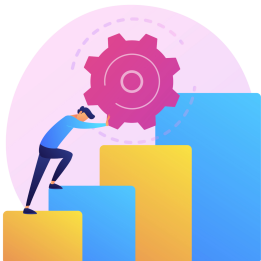
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- TRR $\rightarrow \mathcal{N}$

Exploitable in the Wild?



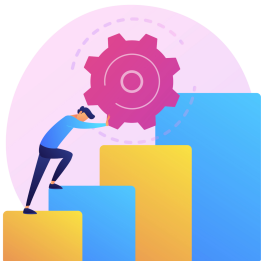
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- **C1**: Allocation of Contiguous Memory
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- **C4**: Bit-Flip Verification



- Mapping from virtual to physical addresses

$$\mathbf{X}_0 = b_8$$

$$\mathbf{X}_1 = b_{12} \oplus b_{16}$$

$$\mathbf{X}_2 = b_{13} \oplus b_{17}$$

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- DRAM addressing function
- Mapping physical address to 16 DRAM banks

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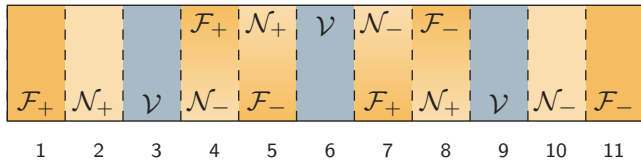
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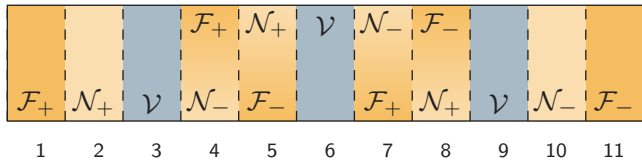
$$\mathbf{X}_3 = b_{14} \oplus b_{18}$$

- Mapping from virtual to physical addresses
- DRAM addressing function
- Mapping physical address to 16 DRAM banks
- **Specific** *bank* access pattern if contiguous memory
- ✓ Extract pattern with a timing side channel

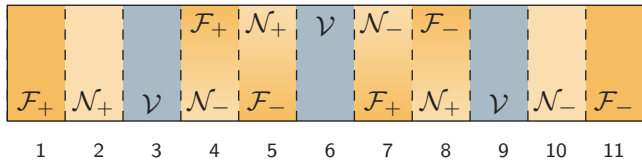




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- Hammer with Half-Double



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- **Spectre** gadget

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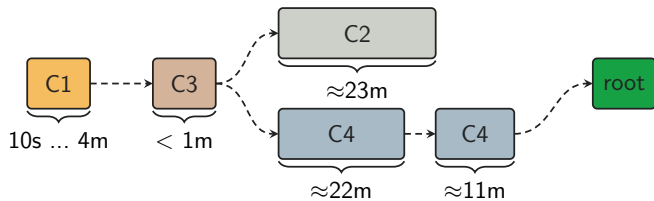
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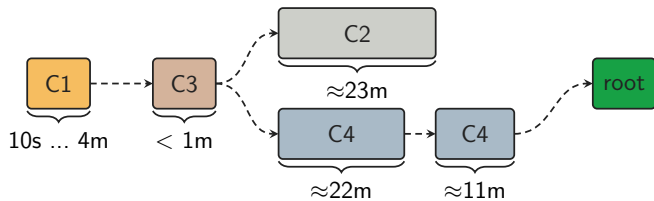
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- **Suppresses** corruption faults

End-to-End Exploit - Timings



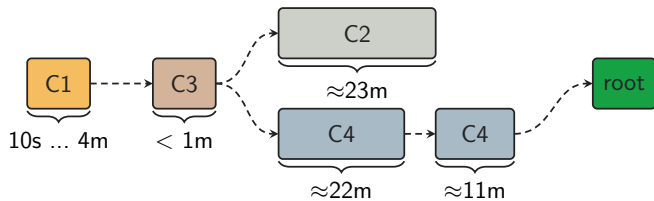
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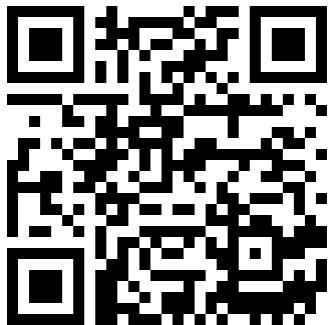
End-to-End Exploit - Timings




- **45** minutes (*Chromebook₂*)
- Full memory read & write primitive
- Deployable **inside** an APP




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- **More details**
 - *Dance*-experiments
 - Contiguous memory Z3 solver
 - Physical address bit recovery
 - ...






- **Open Source**  <https://github.com/IAIK/halfdouble>
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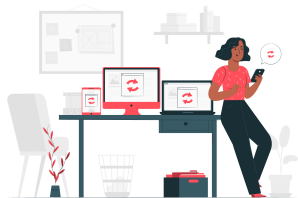


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Read the Paper

-  Pietro Frigo, Emanuele Vannacci, Hasan Hassan, Victor van der Veen, Onur Mutlu, Cristiano Giuffrida, Herbert Bos, and Kaveh Razavi. TRRespass: Exploiting the Many Sides of Target Row Refresh. In: S&P. 2020.
-  Finn de Ridder, Pietro Frigo, Emanuele Vannacci, Herbert Bos, Cristiano Giuffrida, and Kaveh Razavi. SMASH: Synchronized Many-sided Rowhammer Attacks From JavaScript. In: USENIX Security Symposium. 2021.
-  Mark Seaborn and Thomas Dullien. Test DRAM for bit flips caused by the rowhammer problem. Retrieved on July 27, 2015. 2015. URL: <https://github.com/google/rowhammer-test>.

Additional Slides



- Tested **13** DIMMs & devices
- **2** DIMMs affected
 - FPGA analysis
 - Exact numbers
- **5** out of **7** mobile devices affected
 - **Reversed** addressing
 - Unprivileged **flush**
 - Uncachable memory (10x)

Affected Devices - Flip Numbers

System	RAM	$N_{Hammers}$	$UC_{0 \rightarrow 1}$	$UC_{1 \rightarrow 0}$	$Flush_{0 \rightarrow 1}$	$Flush_{1 \rightarrow 0}$
Chromebook ₁	LPDDR4x	23 274	27	40	2	5
Chromebook ₂	LPDDR4x	23 586	235	2379	12	101
OnePlus 5T	LPDDR4x	25 687	2	30	1	24
Pixel 3	LPDDR4x	32 921	11	5	0	0
HTC U11	LPDDR4x	21 840	-	-	3	17